

# Transformation Procedure for Generating Random Variates from the Exponential Distribution

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## **Abstract**

In this paper we introduce a new technique for generating random variates from the exponential distribution. The procedure is based on a proposed mapping that transform an order sample from  $U(0, 1)$  to the  $\lambda \cdot \text{Exp}(1)$ . The method is developed theoretically and assessed practically by the basic Monte Carlo simulation. Comparison is made with that of inverse transform method.

## **Introduction**

In recent years, several procedures are suggested, tested, and used for generating random variates from  $\text{Exp}(\lambda)$ . The procedures of Nayor[1], Toeber[2], Abrams and Dieter[3], and Marsaglia[4] for generating random variates from  $\text{Exp}(\lambda)$  are based on the composition method, Acceptance-rejection method and Fortythe method[5]. These procedures without the benefit of a logarithmic transformation. Neumann's method[6] for generating random variates from  $\text{Exp}(\lambda)$  based on sampling a sequence of i.i.d r.v's from the standard triangular distribution

$$f(x) = \begin{cases} 2x & 0 < x < \lambda \\ \lambda & \lambda < x \leq 1 \\ 2(x-1) & \end{cases}$$

and it is shown that generation of  $n^2$  random exponential variates in such way requires on the average a sequence of 6 random numbers.

Inverse transform method[7] is the one of most common use for generating random variates from  $\text{Exp}(\lambda)$  which can be described by the following algorithm. Note that an exponential variate  $X$  has p.d.f

$$f(x) = \frac{1}{\lambda} e^{-\frac{x}{\lambda}}, 0 < x < \infty$$

$$= 0 \quad , \quad x \leq 0; \quad \lambda > 0$$

By inverse transform method

$$U = F(X) = 1 - e^{-\frac{x}{\lambda}}$$

So that  $X = -\ln(1-U)$ .

Since  $1-U$  is distributed in the same way as  $U$ , we have

$$X = -\ln U$$

## **II. Algorithm**

- 1. Read  $\lambda$ .
- 2. Generating  $U$  from  $U(0,1)$ .
- 3. Set  $X = -\lambda \cdot \ln U$ .

4. Deliver  $X$  as a random variate generated from  $\lambda \cdot \text{Exp}(1)$ .

We note that, if the random variates  $X$  is sampled from  $\text{Exp}(\lambda)$  then the r.v  $Y=\lambda X$  is from  $\text{Exp}(1)$ .

Although the algorithm seem very simple, but the computation of the natural algorithm on a digital computer consist of a power series expansion for the each uniform variates generated.

## **2. Proposition**

Let  $U_1, U_2, \dots, U_n, U_{n+1}, \dots, U_{n+k}$  be a random sample of size  $2n+1$  from  $U(0, 1)$  and let  $W_1, W_2, \dots, W_k$  be the order statistic corresponding to the r.v's  $U_{n+1}, U_{n+2}, \dots, U_{n+k}$ . Assume  $W_0=0$  and  $W_k=1$ , then the r.v's

$$Y_k = -(W_k - W_{k-1}) \ln \prod_{i=1}^k U_i, \quad k = 1, 2, \dots, n$$
 represent

a sample of size  $n$  from  $\text{Exp}(1)$ .

## **Proof**

Let  $X_k = W_k - W_{k-1}$ ,  $k=1, 2, \dots, n-1$  and let  $X_n = \ln \prod_{i=1}^k U_i = \sum_{i=1}^n \ln U_i$ .

The dist. of  $X_n$  can be obtained by using m.g.f technique. Viz  $M_X(t)$  be the m.g.f of  $X$ , then

$$M_{X_n}(t) = E(e^{tX_n}) = E(e^{t \sum_{i=1}^n \ln U_i}) = E\left(\prod_{i=1}^n e^{t \ln U_i}\right) = E\left(\prod_{i=1}^n U_i^t\right) = \prod_{i=1}^n E(U_i^t)$$

because  $U_1, U_2, \dots, U_n$  are independent but for any  $t > U(0, 1)$ ,

$$E(U_i^t) = \int_0^t u^{t-1} du = \frac{1}{1-t}$$

Therefore  $M_{X_n}(t) = \prod_{i=1}^n \frac{1}{1-t} = \frac{1}{(1-t)^n}$  which is the m.g.f of  $C(n, 1)$ , that is  $X_n \sim G(n, 1)$  with p.d.f

$$f_n(x_n) = \frac{1}{\Gamma(n)} x_n^{n-1} e^{-x_n}, \quad 0 < x_n < \infty$$

$$= 0, \quad x_n \geq 0$$

next, let us find the distn. Of r.v's  $X_k - W_k - W_{k+1}$ ,  $k=1,2,\dots,n-1$ .

Since  $U_{i+1}, U_{i+2}, \dots, U_{n+1}$  are independent from  $U(0, 1)$ . Then the joint p.d.f is  $g(u_{i+1}, u_{i+2}, \dots, u_{n+1}) = 1$ ,  $0 < u_i < 1, i=1,2,\dots,n-1$

$$= 0, \quad \text{e.w.}$$

from order statistics theory the joint p.d.f is of  $W_1, W_2, \dots, W_{n-1}$  is

$$h(w_1, w_2, \dots, w_{n-1}) = (n-1)! \cdot 0 < w_1 < w_2 < \dots < w_{n-1} < 1 \\ = 0, \quad \text{e.w.}$$

now the functions  $x_1 - w_1, x_2 - w_2, \dots, x_n - w_{n-1}$  define one-to-one transformation that maps the space  $A = \{(w_1, w_2, \dots, w_{n-1}) : 0 < w_1 < w_2 < \dots < w_{n-1} < 1\}$  on the space  $B = \{(x_1, x_2, \dots, x_{n-1}) : x_i > 0, i=1,2,\dots,n-1\}$

inside  $\sum_{k=1}^{n-1} x_k \leq 1\}$  with inverse transform

$$w_1 = x_1, w_2 = x_1 + x_2, \dots, w_{n-1} = x_1 + x_2 + \dots + x_{n-1}$$

and the jacobian of the transformation

$$J = \frac{\partial(w_1, w_2, \dots, w_{n-1})}{\partial(x_1, x_2, \dots, x_{n-1})} = \begin{vmatrix} 1 & 0 & \dots & 0 \\ 0 & 1 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & 1 \end{vmatrix}_{(n-1) \times (n-1)} = 1$$

Then the joint p.d.f of  $X_1, X_2, \dots, X_n$  is

$$f(x_1, x_2, \dots, x_{n-1}, x_n) = h(x_1, x_1 + x_2, \dots, \sum_{i=1}^{n-1} x_i) J \\ = (n-1)! \cdot x_k \geq 0, k=1,2,\dots,n-1, \quad \sum_{i=1}^{n-1} x_i \leq 1$$

Since the r.v's are independent of  $X_n$ , then the joint p.d.f of

$X_1, X_2, \dots, X_{n-1}$

$$f(x_1, x_2, \dots, x_{n-1}, x_n) = f^*(x_1, x_2, \dots, x_{n-1}) f_n(x_n) \\ = (n-1)! \frac{1}{\Gamma(n)} x_n^{n-1} e^{-x_n} \\ = x_n^{n-1} e^{-x_n}, 0 < x_n < \infty \\ = 0, \quad \text{e.w.}$$

Finally, consider the transformation

$Y_k = X_k / X_n, \quad k=1,2,\dots,n-1$  and  $Y_n = (1-X_1-X_2-\dots-X_{n-1})X_n$

That is

$$Y_1 = X_1 / X_n$$

$$Y_2 = X_2 / X_n$$

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$$Y_{n-1} = X_{n-1} / X_n$$

$$Y_n = (1-X_1-X_2-\dots-X_{n-1})X_n$$

and

$$\sum_{k=1}^n Y_k = (Y_1 + Y_2 + \dots + Y_{n-1})X_n + (1-Y_1-Y_2-\dots-Y_{n-1})X_n = X_n$$

Now, the function  $y_k = x_k / x_n, k=1,2,\dots,n-1$  and

$$y_n = \left( 1 - \sum_{k=1}^{n-1} x_k \right) / x_n$$

define one-to-one transformation that maps the space

$$A = \{(x_1, x_2, \dots, x_{n-1}, x_n) : x_i > 0, i=1,2,\dots,n-1, \sum_{i=1}^{n-1} x_i \leq 1, 0 < x_n < \infty\}$$

onto the space

$$B^* = \{(y_1, y_2, \dots, y_n) : 0 < y_i < \infty, i=1,2,\dots,n\}$$

With inverse

$$X_k = \frac{y_k}{\sum_{i=1}^n y_i}, k=1,2,\dots,n-1, x_n = \sum_{i=1}^n y_i$$

$$J^{-1} = \frac{\partial(y_1, y_2, \dots, y_n)}{\partial(x_1, x_2, \dots, x_n)} = \left( \sum_{i=1}^n y_i \right)^{n-1}$$

Then the joint p.d.f of  $Y_1, Y_2, \dots, Y_n$  is

$$g^*(y_1, y_2, \dots, y_n) = f\left[\frac{y_1}{\sum_{i=1}^n y_i}, \frac{y_2}{\sum_{i=1}^n y_i}, \dots, \frac{y_{n-1}}{\sum_{i=1}^n y_i}, \sum_{i=1}^n y_i\right] \\ = \left( \sum_{i=1}^n y_i \right)^{n-1} e^{\frac{-\sum_{i=1}^n y_i}{\sum_{i=1}^n y_i}} \left( \sum_{i=1}^n y_i \right)^{n-1} \\ = e^{\frac{\sum_{i=1}^n y_i}{\sum_{i=1}^n y_i}}, \quad 0 \leq y_i \leq \infty, \quad i=1,2,\dots,n$$

Which is the joint p.d.f of  $n$  independent r.v's from  $\text{Exp}(1)$ .

Algorithm PT describes the necessary steps for generating r.v's from  $\text{Exp}(1)$  by the proposed procedure

#### PT-Algorithm:

1. Generate  $U_1, U_2, \dots, U_n, U_{n+1}, \dots, U_{n+1}$  from  $U(0, 1)$ .
2. Arrange  $U_{n+1}, \dots, U_{n+1}$  in ascending order of magnitudes by using the order statistics  $W_1, W_2, \dots, W_{n-1}$ .
3. Set  $W_n = 0$  and  $W_{n+1} = 1$ .
4. Put  $Y_k = U_k / \sum_{i=1}^n U_i, k=1,2,\dots,n$ .
5. Deliver  $Y_k, k=1,2,\dots,n$  as a r.v generated from  $\text{Exp}(1)$ .

#### 3. Conclusion:

Comparing PT-Algorithm with PT-Algorithm

1. The advantage of PT-Algorithm, it requires only one computation of  $\ln \prod_{i=1}^n g_i$  for generating  $n$  exponential variates simultaneously while IT-Algorithm requires  $n$  computation of  $\ln g_i$  for each variates  $Y_i$  ( $i=1, 2, \dots, n$ ) separately.
2. The disadvantage of PT-Algorithm, it need  $2n-1$  uniform variates while IT-Algorithm requires only  $n$  uniform variates.
3. PT-Algorithm requires the arrangement of the uniform variates  $U_1 < U_2 < \dots < U_{2n-1}$  to be order statistics  $W_1 < W_2 < \dots < W_{n+1}$  and then calculation  $W_i - W_j$  which also time consuming.
4. Simulating both algorithm, we find that PT-Algorithm is faster than IT-Algorithm for  $n=3, 4, 5, 6$ . The optimal  $n$  is 4.

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